We introduced the basic proof structure of Section 1.2 because truth tables quickly became cumbersome and unrealistic to use. Again, we must introduce more complicated proofs because the large number of definitions and axioms leads to more complex statements which makes the methods of propositional logic unrealistic.

Question 1. Suppose a and b are odd integers. What can be said about their sum a + b? Explain your reasoning.

Question 2. What can be said about the sum of two even integers?

Direct Proofs: The structure of a proof sequence in propositional logic is straightforward: in order to prove $A \Rightarrow C$, we prove a sequence of results.

 $A \Rightarrow B_1 \Rightarrow B_2 \Rightarrow \cdots \Rightarrow B_n \Rightarrow C$

A direct proof in mathematics has the same logic, but we don't usually write such proofs as lists of statement and reasons. Instead, this linear chain of implications is couched in mathematical prose and written in paragraph form. Our first example of a direct proof already appeared in Theorem 1 of Section 1.4.

Example 1. Prove the following statement.

For all real numbers x, if x > 1, then $x^2 > 1$.

Rule of Thumb 1. To prove a statement of the form $(\forall x)(P(x) \rightarrow Q(x))$ directly, begin your proof with a sentence of the form

Let x be [an element of the domain], and suppose P(x).

A direct proof is then a sequence of justified conclusions culminating in Q(x).

Definition 1. An integer x divides an integer y if there is some integer k such that y = kx. We write x|y to denote x divides y. Conversely, we say that y is a multiple of x whenever x divides y.

Axiom 1. If a and b are integers, so are a + b and $a \cdot b$.

Example 2. Prove the following.

For all integers a, b, and c, if a|b and a|c, then a|(b+c).

Example 3. Prove the following.

For all integers a, b, and c, if a|b and b|c, then a|c.

Proof by Contraposition: We have already seen that the contrapositive is equivalent to an implication statement; i.e. $(P(x) \to Q(x)) \iff (\neg Q(x) \to \neg P(x))$. Therefore, using predicate logic, we also have

$$(\forall x)(P(x) \to Q(x)) \iff (\forall x)(\neg Q(x) \to \neg P(x)).$$

Utilization of this fact is called a **proof by contraposition**.

Rule of Thumb 2. To prove a statement of the form $(\forall x)(P(x) \rightarrow Q(x))$ by contraposition, begin your proof with a sentence of the form

Let x be [an element of the domain], and suppose $\neg Q(x)$.

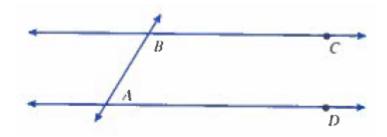
A proof by contraposition is then a sequence of justifies conclusions culminating in $\neg P(x)$.

Example 4. Suppose that x and y are two positive real numbers such that the geometric mean \sqrt{xy} is different from the arithmetic mean $\frac{x+y}{2}$. Show that $x \neq y$.

Theorem 1. The sum of the measures of the angles of any triangle (in Euclidean geometry) is equal to 180° .

Definition 2. Two lines are **parallel** if they do not intersect.

Example 5. Prove that if two lines are cut by a transversal such that a pair of interior angles are supplementary, then the lines are parallel.



Proof by Contradiction: Sometimes even simple-looking statements can be hard to prove directly, with or without contraposition. In this case we can try a **proof by contradiction.** Suppose we want to prove that a statement A is true. Then we argue

$$\neg A \Rightarrow B_1 \Rightarrow B_2 \Rightarrow \cdots \Rightarrow B_n \Rightarrow \mathbf{F},$$

where \mathbf{F} represents a statement that is always false, that is, a contradiction. By taking contrapositives we get the sequence

$$A \Leftarrow \neg B_1 \Leftarrow \neg B_2 \Leftarrow \cdots \Leftarrow \neg B_n \Leftarrow \mathbf{T},$$

where \mathbf{T} is always true, that is, a tautology.

Rule of Thumb 3. To prove a statement A by contradiction, begin your proof with the following sentence:

Suppose, to the contrary, that $\neg A$. You will notice that the negations we practiced in Section 1.3 will come in handy. **Example 6.** In Euclidean geometry, prove that if two lines share a common perpendicular, then the lines are parallel.

Definition 3. An integer n is **even** if n = 2k for some integer k. **Definition 4.** An integer n is **odd** if n = 2k + 1 for some integer k. **Axiom 2.** For all integers n, $\neg(n \text{ is even}) \iff (n \text{ is odd})$.

Lemma 2. Let n be an integer. If n^2 is even, then n is even. Proof.

Example 7. Prove that $\sqrt{2}$ is irrational.

Homework. (Due Oct 1, 2018) Section 1.5: 2, 6, 11, 14, 16

Practice Problems. Section 1.5: 1-9 (odd), 13, 17-20, 25